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PRIMALITY TEST WITH PAIR OF LUCAS SEQUENCES

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Abstract. Lucas sequences and their applications play vital role in the study of primality tests in number theory.

There are several known tests for primality of positive integer N using Lucas sequences which are based on fac-

torization of $(N \pm 1)$ [2] [13]. In this paper we give a primality test for odd positive integer N > 1 by using the

set $L(\Delta, N)$ where $L(\Delta, N)$ is a set of S(N) distinct pair of Lucas sequences $(V_n(a, 1), U_n(a, 1))$, where S(N) for

 $N = p_1^{e_1} \cdot p_2^{e_2} \dots p_s^{e_s}$ is given as $S(N) = \text{LCM}\left[\left\{p_i^{e_i-1}\left(p_i - \left(\frac{\Delta}{p_i}\right)\right)\right\}_{i=1}^s\right]$ and $\Delta = a^2 - 4$ for some fixed integer a.

Keywords: Lucas sequences; primality testing.

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1. Introduction

Lucas sequences are recurrence relations. Many studies on properties of Lucas sequences,

and their connections with topics like trigonometric functions, Chebyshev's functions, Dickson

functions, continued fractions are known [2][7]. The primality test for a positive integer N using

Lucas sequences was first initiated by Lucas and later developed further by Lehmer, was based

on factorization of $(N \pm 1)$ [2][8]. In this paper we give a primality test for odd positive integer

N > 1 by using the set $L(\Delta, N)$ where $L(\Delta, N)$ is a set of S(N) distinct pair of Lucas sequences

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 $\left(V_n(a,1), U_n(a,1)\right)$, where S(N) for $N=p_1^{e_1}.p_2^{e_2}...p_s^{e_s}$ is given as $S(N)=\mathrm{LCM}\left[\left\{p_i^{e_i-1}\left(p_i-\left(\frac{\Delta}{p_i}\right)\right)\right\}_{i=1}^s\right]$ and $\Delta=a^2-4$ for some fixed integer a. In the following we describe the pair of Lucas sequences and their properties.

Definition 1.1. Let a and b be two integers, α a root of the polynomial $x^2 - ax + b$ in $\mathbb{Q}(\sqrt{\Delta})$ for $\Delta = a^2 - 4b$ a non square, writing $\alpha = \frac{a + \sqrt{\Delta}}{2}$ and its conjugate $\beta = \frac{a - \sqrt{\Delta}}{2}$ we have $\alpha + \beta = a$, $\alpha\beta = b$, $\alpha - \beta = \sqrt{\Delta}$, and the Lucas sequences $V_n(a,b)$ and $U_n(a,b)$, $n \ge 0$ are defined as

$$egin{cases} V_n(a,b) = \pmb{lpha}^n + \pmb{eta}^n, \ U_n(a,b) = rac{\pmb{lpha}^n - \pmb{eta}^n}{\pmb{lpha} - \pmb{eta}}. \end{cases}$$

In particular, $V_0(a,b) = 2$, $V_1(a,b) = a$ and $U_0(a,b) = 0$, $U_1(a,b) = 1$.

 $V_n(a,b)$ and $U_n(a,b)$ are given by following recurrence sequences:

$$\begin{cases} V_n(a,b) = aV_{n-1}(a,b) - bV_{n-2}(a,b), \\ U_n(a,b) = aU_{n-1}(a,b) - bU_{n-2}(a,b). \end{cases}$$

Lucas sequences satisfy the following properties [3] [4] [9]:

$$(1) \left(V_{2n}(a,b), U_{2n}(a,b) \right) = \left((V_n(a,b))^2 - 2b^n, U_n(a,b)V_n(a,b) \right).$$

(2)
$$(V_n^2(a,b), U_n^2(a,b))' = (\Delta(U_n(a,b))^2 + 4b^n, U_{n-1}(a,b)U_{n+1}(a,b) + b^{n-1}).$$

(3)
$$(2V_{m+n}(a,b), 2U_{m+n}(a,b)) = (V_m(a,b)V_n(a,b) + \Delta U_m(a,b)U_n(a,b),$$

 $U_m(a,b)V_n(a,b) + U_n(a,b)V_m(a,b)) \, \forall \, m \geq n.$

$$(4) \left(V_{m+n}(a,b), U_{m+n}(a,b) \right) = \left(V_m(a,b)V_n(a,b) - b^n V_{m-n}(a,b), \\ U_m(a,b)V_n(a,b) - b^n U_{m-n}(a,b) \right), \ \forall \ m \ge n.$$

In particular for b = 1 the above properties can be written as

(1)
$$\left(V_{2n}(a,1), U_{2n}(a,1)\right) = \left(\left(V_n(a,1)\right)^2 - 2, U_n(a,1)V_n(a,1)\right).$$

(2)
$$(V_n^2(a,1), U_n^2(a,1)) = (\Delta(U_n(a,1))^2 + 4, U_{n-1}(a,1)U_{n+1}(a,1) + 1).$$

(3)
$$\left(2V_{m+n}(a,1), 2U_{m+n}(a,1)\right) = \left(V_m(a,1)V_n(a,1) + \Delta U_m(a,1)U_n(a,1), \right)$$

$$U_m(a,1)V_n(a,1)+U_n(a,1)V_m(a,1)$$
 $\forall m \geq n$.

(4)
$$\left(V_{m+n}(a,1), U_{m+n}(a,1)\right) = \left(V_m(a,1)V_n(a,1) - V_{m-n}(a,1), U_m(a,1)V_n(a,1) - U_{m-n}(a,1)\right), \forall m \ge n.$$

Definition 1.2. If $N = p_1^{e_1} . p_2^{e_2} ... p_s^{e_s}$ and $\Delta = a^2 - 4b$ for some fixed integer a such that $(N,\Delta) = 1$ then we have define $S(N) = \text{LCM}[n_1, n_2, \dots, n_s]$ where $n_i = p_i^{e_i - 1} \left(p_i - \left(\frac{\Delta}{p_i} \right) \right)$ for all $1 \le i \le s$ and $\left(\frac{\Delta}{p_i}\right)$ is the Legendre symbol [12] of Δ with respect to the prime p_i .

$$(1) \left(V_{S(N)}(a,b), U_{S(N)}(a,b)\right) \equiv \left(2b^{\frac{k(1-\varepsilon)}{2}}, 0\right) \mod N.$$

$$(2) \left(V_{S(N)t}(a,b), U_{S(N)t}(a,b)\right) \equiv \left(2b^{\frac{k(1-\varepsilon)}{2}}, 0\right) \mod N.$$

In particular for b = 1, we have

(1)
$$\left(V_{S(N)}(a,1), U_{S(N)}(a,1)\right) \equiv (2,0) \mod N.$$

(2) $\left(V_{S(N)t}(a,1), U_{S(N)t}(a,1)\right) \equiv (2,0) \mod N.$

(2)
$$\left(V_{S(N)t}(a,1), U_{S(N)t}(a,1)\right) \equiv (2,0) \mod N$$

In the following an algorithm is given for computation of Lucas sequences $(V_n(a,1),U_n(a,1))$ using Lucas addition chain as in [11]. This algorithm gives Lu-

Algorithm 1 Evaluate $(V_n(a,1), U_n(a,1))$

step 0: (Initialize) Set
$$N \leftarrow \frac{n}{2^{k-i}}$$
 where $k = \lfloor \log n \rfloor, i = 0, 1, 2, \dots, k \mid Y \leftarrow 1, Z \leftarrow 2$

step 1: (Value N) $N \leftarrow \frac{n}{2^{k-i}}$ and determine whether N is even or odd, if N is even skip to step 4.

step 2: set
$$Y \leftarrow 2Y + 1$$
 and $Z \leftarrow 2Z$

step 3: [N = n], if N = n the algorithm terminates with Y as the answer.

step 4: set
$$Y \leftarrow 2Y, Z \leftarrow Y + 1$$
 and return to step 1.

step 5: [initialize
$$(V_n(a,1), U_n(a,1)]$$
 set $V_0(a,1) = 2, V_1(a,1) = a$ and $U_0(a,1) = 0, U_1(a,1) = 1$

step 6: For *i* from 0 to *k* set
$$n \leftarrow y + z$$

compute $V_{y+z}(a,1) \leftarrow V_y(a,1)V_z(a,1) - V_{y-z}(a,1)$

and
$$U_{v+z}(a,1) \leftarrow U_v(a,1)V_z(a,1) - U_{v-z}(a,1)$$

cas addition chain [5] [11] $\{e_{-1}, e_0, e_1, e_1 + 1, \dots e_{k-1} - 1, e_{k-1}, e_k\}$ and evaluates $\{(V_{e_{-1}}, U_{e_{-1}}), (V_{e_0}, U_{e_0}), (V_{e_1}, U_{e_1}), \dots (V_{e_{t-1}}, U_{e_{t-1}}, (V_{e_t}, U_{e_t}))\}$ for all $t = 0, 1, \dots, k$ by using the formulas $V_{v+z}(a,1)$ and $U_{v+z}(a,1)$

2. PRIMALITY OF N WITH S(N)

In the following we prove a theorem on primality of N with S(N) [7] [9].

Theorem 2.1. If *N* is odd positive integer then *N* is prime if and only if $S(N) = N - (\frac{\Delta}{N})$ for all Δ with $(\Delta, N) = 1$.

Proof. Let N be a prime number then by definition for N=p, and any Δ with $(\Delta,N)=1$ we have $S(N)=S(p)=p-\left(\frac{\Delta}{p}\right)=N-\left(\frac{\Delta}{N}\right)$. Conversely suppose $S(N)=N-\left(\frac{\Delta}{N}\right)$, now if N is composite then for $N=\prod_{i=1}^s p_i^{e_i}$ we have the two cases (i) s=1 with $e_1\geq 2$ and (ii) $s\geq 2$. In case (i) for $s=1,e_1\geq 2$, we have $N=p_1^{e_1},e_1\geq 2$ and $S(N)=S(p_1^{e_1})=p_1^{e_1-1}\left(p-\left(\frac{\Delta}{p}\right)\right)$, therefore as $e_1\geq 2$ we have $p_1|S(N)$ but note $p_1\nmid p_1^{e_1}\pm 1$ i.e. $p_1\nmid N-\left(\frac{\Delta}{N}\right)$ a contradiction to $S(N)=N-\left(\frac{\Delta}{N}\right)$, hence N is not in case (i). Now if N is as in case (ii) we have $N=\prod_{i=1}^s p_i^{e_i}$, with $s\geq 2$ and $S(N)=\mathrm{LCM}\left[\left\{p_i^{e_i-1}\left(p_i-\left(\frac{\Delta}{p_i}\right)\right)\right\}_{i=1}^s\right]$. Now as p_i' s are odd, $(\Delta,N)=1$ and $\left(p_i-\left(\frac{\Delta}{p_i}\right)\right)$ are even, we note in the following that S(N)< N-1:

$$S(N) = LCM \left[\left\{ p_i^{e_i - 1} \left(p_i - \left(\frac{\Delta}{p_i} \right) \right) \right\}_{i=1}^{s} \right]$$

$$= 2 LCM \left[\left\{ p_i^{e_i - 1} \left(p_i - \left(\frac{\Delta}{p_i} \right) \right) \right\}_{i=1}^{s} \right]$$

$$\leq 2 \prod_{i=1}^{s} \left[\frac{1 - \frac{1}{p_i} \left(\frac{\Delta}{p_i} \right)}{2} p_i^{e_i} \right]$$

$$= 2N \prod_{i=1}^{s} \frac{1}{2} \left[1 + \frac{1}{p_i} \right]$$

$$\leq 2N \frac{1}{2^s} \prod_{i=1}^{s} \left[1 + \sum_{i} \frac{1}{p_i} + \sum_{i,j} \frac{1}{p_i p_j} + \dots + \sum_{i,j,\dots,s} \frac{1}{p_i p_j \dots p_s} \right]$$

$$\leq 2N \frac{1}{2^s} \left[1 + \sum_{i} \frac{1}{5} + \sum_{i,j} \frac{1}{5^2} + \dots + \sum_{i,j,\dots,s} \frac{1}{5^s} \right] \text{ as } p_i \geq 5$$

$$= 2N \frac{1}{2^s} \left[1 + s_{c_1} \left(\frac{1}{5} \right) + s_{c_2} \left(\frac{1}{5^2} \right) + \dots + s_{c_s} \left(\frac{1}{5^s} \right) \right]$$

$$= 2N \frac{1}{2^s} \left(1 + \frac{1}{5} \right)^s$$

$$= 2N \left(\frac{3}{5} \right)^s$$

$$\leq 2N \left(\frac{3}{5} \right)^2 \text{ as } s \geq 2$$

$$< 2N \left(\frac{2}{5} \right) = \frac{4N}{5} < N - 1$$

Therefore S(N) < N-1, and as (N-1) < (N+1) we also have S(N) < N+1 in particular we have $S(N) < N-(\frac{\Delta}{N})$ which is a contradiction to $S(N) = N-(\frac{\Delta}{N})$, therefore N is not in case (ii) as well, therefore N is not composite. Hence N is prime.

3. PRIMALITY TEST WITH PAIR OF LUCAS SEQUENCES

Notation 3.1. Let N be a positive integer and $\Delta = a^2 - 4$ for some positive integer a such that $(N, \Delta) = 1$, then the set of all the pairs of Lucas sequences is denoted as $L(\Delta, N)$ and is given as $L(\Delta, N) = \left\{ \left(V_n(a, 1), U_n(a, 1) \right) : 1 \le n \le S(N) \right\}.$

The following theorem assures that all the pairs in $L(\Delta, N)$ are distinct modulo N and $|L(\Delta, N)| = S(N)$.

Theorem 3.2. $(V_r(a,1), U_r(a,1)) \equiv (V_0(a,1), U_0(a,1)) \mod N$ if and only if $r \equiv 0 \mod S(N)$ [10].

Proof. Suppose $r \equiv 0 \mod S(N)$, then we have r = S(N)t, for some integer t and $\Big(V_r(a,1),U_r(a,1)\Big) \equiv \Big(V_{S(N)t}(a,1),U_{S(N)t}(a,1)\Big) \equiv \Big(V_0(a,1),U_0(a,1)\Big) \mod N$, therefore $r \equiv 0 \mod S(N)$ implies $\Big(V_r(a,1),U_r(a,1)\Big) \equiv \Big(V_0(a,1),U_0(a,1)\Big) \mod N$. Conversely suppose $\Big(V_r(a,1),U_r(a,1)\Big) \equiv \Big(V_0(a,1),U_0(a,1)\Big) \mod N$, then for $N = p_1^{e_1}.p_2^{e_2}...p_s^{e_s}$ note $\Big(V_r(a,1),U_r(a,1)\Big) \equiv \Big(V_0(a,1),U_0(a,1)\Big) \mod p_i^{e_i}$ for all $i = 1,2,\ldots,s$. We now show in the following that $p_i^{e_i-1}\Big(p_i-\Big(\frac{\Delta}{p_i}\Big)\Big)$ divides r for all $i = 1,2,\ldots,s$. For all Δ with $(\Delta,N)=1$ as $(\Delta,p_i^{e_i})=1$ using Euler's criterion we have

$$\alpha^{p_i} \equiv \begin{cases} \alpha \mod p_i & \text{if } \left(\frac{\Delta}{p_i}\right) = 1, \\ \beta \mod p_i & \text{if } \left(\frac{\Delta}{p_i}\right) = -1 \end{cases}$$

$$\Rightarrow \alpha^{p_i} \equiv \begin{cases} \alpha + kp_i & \text{if } \left(\frac{\Delta}{p_i}\right) = 1, \\ \beta + kp_i & \text{if } \left(\frac{\Delta}{p_i}\right) = -1 \end{cases}$$

Therefore for i = 1, 2, ..., s we have,

$$\alpha^{p_i^{e_i}} = (\alpha^{p_i})^{p_i^{e_i-1}} = (\alpha + kp_i)^{p_i^{e_i-1}}$$

$$= \alpha^{p_i^{e_i-1}} + \binom{p_i^{e_i-1}}{1} \alpha^{p_i^{e_i-1}-1} (kp_i) + \binom{p_i^{e_i-1}}{2} \alpha^{p_i^{e_i-1}-2} (kp_i)^2 + \dots + \binom{p_i^{e_i-1}}{p_i^{e_i-1}-1} \alpha(kp_i)^{p_i^{e_i-1}-1} + (kp_i)^{p_i^{e_i-1}}$$

$$\Rightarrow \alpha^{p_i^{e_i}} \equiv \alpha^{p_i^{e_i-1}} \mod p_i^{e_i} \text{ if } \left(\frac{\Delta}{p_i}\right) = 1.$$

Similarly,

$$egin{aligned} lpha^{p_i^{e_i}} &= (lpha^{p_i})^{p_i^{e_i-1}} \ &= (eta + kp_i)^{p_i^{e_i-1}} \ &\Rightarrow lpha^{p_i^{e_i}} \equiv eta^{p_i^{e_i-1}} \mod p_i^{e_i} ext{ if } \left(rac{\Delta}{p_i}
ight) = 1 \end{aligned}$$

therefore,

$$lpha^{p_i^{e_i}} \equiv egin{cases} lpha^{p_i^{e_i-1}} & mod p_i^{e_i} & mod rac{\Delta}{p_i} = 1, \ eta^{p_i^{e_i-1}} & mod p_i^{e_i} & mod rac{\Delta}{p_i} = -1 \end{cases}$$

 $\begin{array}{l} \text{now } \alpha^{p_i^{e_i}} \equiv \alpha^{p_i^{e_i-1}} \mod p_i^{e_i} \text{ if } \left(\frac{\Delta}{p_i}\right) = 1 \Rightarrow \alpha^{p_i^{e_i}(p-1)} \equiv 1 \mod p_i^{e_i} \text{ if } \left(\frac{\Delta}{p_i}\right) = 1 \\ \text{and } \alpha^{p_i^{e_i}} \equiv \beta^{p_i^{e_i-1}} \mod p_i^{e_i} \text{ if } \left(\frac{\Delta}{p_i}\right) = -1 \Rightarrow \alpha^{p_i^{e_i}(p+1)} \equiv 1 \mod p_i^{e_i} \text{ if } \left(\frac{\Delta}{p_i}\right) = -1 \\ \text{therefore note } p_i^{e_i} \text{ is smallest such that} \end{array}$

$$lpha^{p_i^{e_i}} \equiv egin{cases} lpha^{p_i^{e_i-1}} \mod p_i^{e_i} ext{ if } \left(rac{\Delta}{p_i}
ight) = 1, \ eta^{p_i^{e_i-1}} \mod p_i^{e_i} ext{ if } \left(rac{\Delta}{p_i}
ight) = -1 \end{cases}$$

 $\Rightarrow p_i^{e_i} \text{ is smallest such that } \alpha^{p_i^{e_i-1}\left(p_i-\left(\frac{\Delta}{p_i}\right)\right)} \equiv 1 \mod p_i^{e_i}.$ Now note $\left(V_r(a,1),U_r(a,1)\right) \equiv \left(V_0(a,1),U_0(a,1)\right) \mod p_i^{e_i}$ $\Rightarrow V_r(a,1) \equiv V_0(a,1) \mod p_i^{e_i} \text{ and } U_r(a,1) \equiv U_0(a,1) \mod p_i^{e_i}$ $\Rightarrow \alpha^r + \beta^r \equiv 2 \mod p_i^{e_i} \text{ and } \frac{\alpha^r - \beta^r}{\alpha - \beta} \equiv 0 \mod p_i^{e_i}$ $\Rightarrow \alpha^r + \beta^r \equiv 2 \mod p_i^{e_i} \text{ and } \alpha^r - \equiv \beta^r \mod p_i^{e_i}$ $\Rightarrow 2\alpha^r \equiv 2 \mod p_i^{e_i}$

 $\Rightarrow \alpha^r \equiv 1 \mod p_i^{e_i}$, therefore $p_i^{e_i-1} \left(p_i - \left(\frac{\Delta}{p_i} \right) \right)$ divides r for $i = 1, 2, \dots, s$ therefore r is a

common multiple of $p_i^{e_i-1}\Big(p_i-\Big(\frac{\Delta}{p_i}\Big)\Big)$ and as S(N) is the $\operatorname{LCM}\Big[\Big\{p_i^{e_i-1}\Big(p_i-\Big(\frac{\Delta}{p_i}\Big)\Big)\Big\}_{i=1}^s\Big]$ we have S(N)|r which implies $r\equiv 0 \mod S(N)$, therefore we have $\Big(V_r(a,1),U_r(a,1)\Big)\equiv \Big(V_0(a,1),U_0(a,1)\Big)\mod N$ implies $r\equiv 0\mod S(N)$.

Now in the following theorem we propose a primality test for N by using the pair of Lucas sequences.

Theorem 3.3. For any odd positive integer N > 1, N is prime if and only if $\left(V_{N-\left(\frac{\Delta}{N}\right)}(a,1), U_{N-\left(\frac{\Delta}{N}\right)}(a,1)\right) \equiv \left(V_0(a,1), U_0(a,1)\right) \mod N$, for all Δ with $(\Delta,N) = 1$.

Proof. Let N be prime number, then by definition $S(N) = N - \left(\frac{\Delta}{N}\right)$ for all Δ with $(\Delta, N) = 1$ and as $S(N) \equiv 0 \mod S(N)$, by Theorem 3.1 we have $\left(V_{S(N)}(a,1), U_{S(N)}(a,1)\right) \equiv \left(V_0(a,1), U_0(a,1)\right) \mod N$ $\Rightarrow \left(V_{N-\left(\frac{\Delta}{N}\right)}(a,1), U_{N-\left(\frac{\Delta}{N}\right)}(a,1)\right) \equiv \left(V_0(a,1), U_0(a,1)\right) \mod N$ $\therefore N$ is prime implies $\left(V_{N-\left(\frac{\Delta}{N}\right)}(a,1), U_{N-\left(\frac{\Delta}{N}\right)}(a,1), U_{N-\left(\frac{\Delta}{N}\right)}(a,1)\right) \equiv \left(V_0(a,1), U_0(a,1)\right) \mod N$, for all Δ with $(\Delta, N) = 1$.

Conversely let $\left(V_{N-\left(\frac{\Delta}{N}\right)}(a,1),U_{N-\left(\frac{\Delta}{N}\right)}(a,1)\right)\equiv \left(V_{0}(a,1),U_{0}(a,1)\right)\mod N$, for all Δ with $(\Delta,N)=1$ then by Theorem 3.1 we have $N-\left(\frac{\Delta}{N}\right)\equiv 0\mod S(N)$, therefore we have $S(N)\mid N-\left(\frac{\Delta}{N}\right)$, for all Δ with $(N,\Delta)=1$.

If possible suppose N is not a prime, then we have the cases that N is composite and not square-free or N is composite and squarefree. First suppose N is composite and not squarefree then $N=p_1^{e_1}.p_2^{e_2}\dots p_s^{e_s}$, for p_1,p_2,\dots,p_s are distinct primes and $s\geq 1$ with $e_i>1$ for some $1\leq i\leq s$, that is $e_i-1>0$ then for some $1\leq i\leq s$, therefore $S(N)=\mathrm{LCM}[p_1^{e_1-1}(p_1-\left(\frac{\Delta}{p_1}\right)),p_2^{e_2-1}(p_2-\left(\frac{\Delta}{p_2}\right),\dots,p_s^{e_s-1}(p_s-\left(\frac{\Delta}{p_s}\right))]\Rightarrow p_i^{e_i-1}\nmid S(N)$ for some $1\leq i\leq s$. Further for Δ with $(\Delta,N)=1$, we have $N-\left(\frac{\Delta}{N}\right)=N\pm 1$ and $p_i\nmid N$ for all $1\leq i\leq s$ note $p_i\nmid (N-\left(\frac{\Delta}{N}\right))$ for all $1\leq i\leq s$. Therefore as $p_i\mid S(N)$ for some $1\leq i\leq s$ and $p_i\nmid (N-\left(\frac{\Delta}{N}\right))$ for all $1\leq i\leq s$ note $S(N)\nmid N-\left(\frac{\Delta}{N}\right)$ which is contradiction to $S(N)\mid N-\left(\frac{\Delta}{N}\right)$ for all Δ with $(N,\Delta)=1$, therefore N is composite and not squarefree is not possible. Now if N is composite and squarefree then $N=p_1.p_2\dots p_s$ for $p_i's$ are distinct primes and s>1, therefore writing $N=p_ip_jt$ for $p_i\neq p_j$ for some $1\leq i,j\leq s$ and for $p_i>3$ we have

$$P_i - \left(\frac{\Delta}{p_i}\right) \equiv 0 \mod p_i - \left(\frac{\Delta}{p_i}\right)$$

$$\Rightarrow p_i \equiv \left(\frac{\Delta}{p_i}\right) \mod p_i - \left(\frac{\Delta}{p_i}\right)$$

$$\Rightarrow N = p_{i}p_{j}t \equiv p_{j}t \left(\frac{\Delta}{p_{i}}\right) \mod p_{i} - \left(\frac{\Delta}{p_{i}}\right)$$

$$\Rightarrow N \equiv p_{j}t \left(\frac{\Delta}{p_{i}}\right) \mod p_{i} - \left(\frac{\Delta}{p_{i}}\right)$$

$$\Rightarrow N - \left(\frac{\Delta}{p_{i}p_{j}t}\right) \equiv p_{j}t \left(\frac{\Delta}{p_{i}}\right) - \left(\frac{\Delta}{p_{i}p_{j}t}\right) \mod p_{i} - \left(\frac{\Delta}{p_{i}}\right)$$

$$\Rightarrow N - \left(\frac{\Delta}{N}\right) \equiv p_{j}t \left(\frac{\Delta}{p_{i}}\right) - \left(\frac{\Delta}{p_{i}}\right) \left(\frac{\Delta}{p_{j}t}\right) \mod p_{i} - \left(\frac{\Delta}{p_{i}}\right)$$

$$\Rightarrow N - \left(\frac{\Delta}{N}\right) \equiv \left(\frac{\Delta}{p_{i}}\right) \left(p_{j}t - \left(\frac{\Delta}{p_{j}t}\right)\right) \mod p_{i} - \left(\frac{\Delta}{p_{i}}\right)$$

$$\Rightarrow p_{j}t - \left(\frac{\Delta}{p_{j}t}\right) \equiv 0 \mod p_{i} - \left(\frac{\Delta}{p_{i}}\right) \operatorname{as} p_{i} - \left(\frac{\Delta}{p_{i}}\right) | N - \left(\frac{\Delta}{N}\right)$$

Therefore we have $p_i - \left(\frac{\Delta}{p_i}\right) \mid p_j t - \left(\frac{\Delta}{p_j t}\right)$, for all Δ with $(\Delta, N) = 1$, but note this is a contradiction as there are some Δ such that $p_i - \left(\frac{\Delta}{p_i}\right) \nmid p_j t - \left(\frac{\Delta}{p_j t}\right)$ which is seen in the following:

For a_1, a_2 with $a_1 \equiv a_2 \mod p_i$ we have for $\Delta_1 = a_1^2 - 4$, $\Delta_2 = a_2^2 - 4$ such that $\left(\frac{\Delta_1}{N}\right) \neq \left(\frac{\Delta_2}{N}\right)$, then $\left(\frac{\Delta_1}{p_i}\right) = \left(\frac{\Delta_2}{p_i}\right)$ and $\left(\frac{\Delta_1}{p_jt}\right) \neq \left(\frac{\Delta_2}{p_jt}\right)$ and we have the following cases;

(i)
$$\left(\frac{\Delta_1}{p_i}\right) = \left(\frac{\Delta_2}{p_i}\right) = 1$$
 and $\left(\frac{\Delta_1}{p_j t}\right) = 1, \left(\frac{\Delta_2}{p_j t}\right) = -1.$

(ii)
$$\left(\frac{\Delta_1}{p_i}\right) = \left(\frac{\Delta_2}{p_i}\right) = 1$$
 and $\left(\frac{\Delta_1}{p_j t}\right) = -1, \left(\frac{\Delta_2}{p_j t}\right) = 1.$

(iii)
$$\left(\frac{\Delta_1}{p_i}\right) = \left(\frac{\Delta_2}{p_i}\right) = -1$$
 and $\left(\frac{\Delta_1}{p_j t}\right) = 1, \left(\frac{\Delta_2}{p_j t}\right) = -1.$

(iv)
$$\left(\frac{\Delta_1}{p_i}\right) = \left(\frac{\Delta_2}{p_i}\right) = -1$$
 and $\left(\frac{\Delta_1}{p_j t}\right) = -1$, $\left(\frac{\Delta_2}{p_j t}\right) = 1$.

in all the cases note either $p_i - \left(\frac{\Delta_1}{p_i}\right) \nmid p_j t - \left(\frac{\Delta_1}{p_j t}\right)$ or $p_i - \left(\frac{\Delta_2}{p_i}\right) \mid p_j t - \left(\frac{\Delta_2}{p_j t}\right)$

i.e.
$$(p_i - 1) \nmid (p_j t - 1)$$
 or $(p_i - 1) \nmid (p_j t + 1)$

$$(p_i-1) \nmid (p_jt+1) \text{ or } (p_i-1) \nmid (p_jt-1)$$

$$(p_i+1) \nmid (p_jt-1) \text{ or } (p_i+1) \nmid (p_jt+1)$$

$$(p_i + 1) \nmid (p_i t + 1) \text{ or } (p_i + 1) \nmid (p_i t - 1)$$

respectively as $(p_i - 1) \mid (p_j t - 1)$ and $(p_i - 1) \nmid (p_j t + 1)$ implies $(p_i - 1) \mid \pm 2$ which implies $p_i = 1$ or 3, a contradiction which implies N is not a composite and squarefree number, therefore N is prime.

In the following, an algorithm is given for evaluating Lucas sequences $(V_n(a,1), U_n(a,1))$ and test for primality of N.

Algorithm 2 Evaluate $(V_n(a,1), U_n(a,1))$ and test for primality of N

step 0: (Initialize) Set
$$N \leftarrow \frac{n}{2^{k-i}}$$
 where $k = \lfloor \log n \rfloor, i = 0, 1, 2, ..., k$
 $Y \leftarrow 1, Z \leftarrow 2$

step 1: (Value N) $N \leftarrow \frac{n}{2^{k-i}}$ and determine whether N is even or odd, if N is even skip to step 4.

step 2: set
$$Y \leftarrow 2Y + 1$$
 and $Z \leftarrow 2Z$

step 3: [N = n], if N = n the algorithm terminates with Y as the answer.

step 4: set
$$Y \leftarrow 2Y$$
, $Z \leftarrow Y + 1$ and return to step 1.

step 5: [initialize
$$(V_n(a,1), U_n(a,1)]$$
 set $V_0(a,1) = 2, V_1(a,1) = a$ and $U_0(a,1) = 0, U_1(a,1) = 1$

step 6: For
$$i$$
 from 0 to k set $n \leftarrow x + y$
compute $V_{y+z}(a,1) \leftarrow V_y(a,1)V_z(a,1) - V_{y-z}(a,1)$
and $U_{y+z}(a,1) \leftarrow U_y(a,1)V_z(a,1) - U_{y-z}(a,1)$

step 7: compute $\left(V_{N\pm 1}(a,1), U_{N\pm 1}(a,1)\right) \mod N$, if it is $\left(V_0(a,1), U_0(a,1)\right) \mod N$ then N is prime otherwise N is composite.

Example 3.4. Let N = 2883155, a = 41 then $\Delta = 1677$ such that $\left(\frac{\Delta}{N}\right) = \left(\frac{1677}{2883155}\right) = -1$ Now compute $\left(V_{N-\left(\frac{\Delta}{N}\right)}(a,1), U_{N-\left(\frac{\Delta}{N}\right)}(a,1)\right) \mod N$ $\equiv \left(V_{2883155+1}(41,1), U_{2883155+1}(41,1)\right) \mod 2883155$ $\equiv \left(V_{2883156}(41,1), U_{2883156}(41,1)\right) \mod 2883155$ $\equiv \left(V_{276}(41,1), U_{276}(41,1)\right) \mod 2883155$ $\equiv \left(80, 192239\right) \mod 2883155$ so, $\left(V_{N-\left(\frac{\Delta}{N}\right)}(a,1), U_{N-\left(\frac{\Delta}{N}\right)}(a,1)\right) \not\equiv \left(V_{0}(a,1), U_{0}(a,1)\right) \mod N$, therefore N is not a prime.

Example 3.5. Let
$$N=104701$$
, $a=64$ then $\Delta=4092$ such that $\left(\frac{\Delta}{N}\right)=\left(\frac{4092}{104701}\right)=1$
Now compute $\left(V_{N-\left(\frac{\Delta}{N}\right)}(a,1),U_{N-\left(\frac{\Delta}{N}\right)}(a,1)\right)\mod N$
 $\equiv \left(V_{104701-1}(64,1),U_{104701-1}(64,1)\right)\mod 104701$
 $\equiv \left(V_{104700}(64,1),U_{104700}(64,1)\right)\mod 104701$
 $\equiv (2,0)\mod 2883155$
so, $\left(V_{N-\left(\frac{\Delta}{N}\right)}(64,1),U_{N-\left(\frac{\Delta}{N}\right)}(64,1)\right)\equiv \left(V_{0}(64,1),U_{0}(64,1)\right)\mod N$, therefore N is a prime.

Note 3.6. The primality test is independent of choice of a and Δ .

Example 3.7. List of $L(\Delta, N)$ for $\left(\frac{\Delta}{N}\right) = 1$ and $\left(\frac{\Delta}{N}\right) = -1$ for composite and prime N with respect to S(N) is given in the following tables depicting that the primality test for N = 33 and 37 is independent of a and Δ .

$L(\Delta, 33)$				
$a = 12, \Delta = 8, (\frac{\Delta}{N}) =$		$a=18, \Delta=23, (\frac{\Delta}{N})=$	= -1 and S(N) = 20	
(V_0, U_0)	(2,0)	(V_0, U_0)	(2,0)	
(V_1, U_1)	(12,1)	(V_1, U_1)	(18,1)	
(V_2, U_2)	(10,12)	(V_2, U_2)	(25,18)	
(V_3, U_3)	(9,11)	(V_3, U_3)	(3,26)	
(V_4, U_4)	(32,21)	(V_4, U_4)	(29,21)	
(V_5, U_5)	(12,10)	(V_5, U_5)	(24,22)	
(V_6, U_6)	(13,0)	(V_6, U_6)	(7,12)	
(V_7, U_7)	(12,23)	(V_7, U_7)	(3,29)	
(V_8, U_8)	(32,12)	(V_8, U_8)	(14,15)	
(V_9, U_9)	(9,22)	(V_9, U_9)	(18,10)	
(V_{10}, U_{10})	(10,21)	(V_{10}, U_{10})	(13,0)	
(V_{11}, U_{11})	(12,32)	(V_{11}, U_{11})	(19,34)	
(V_{12}, U_{12})	(2,0)	(V_{12}, U_{12})	(14,18)	
(V_{13}, U_{13})	(12,1)	(V_{13}, U_{13})	(3,4)	
(V_{14}, U_{14})	(10,12)	(V_{14}, U_{14})	(7,21)	
(V_{15}, U_{15})	(9,11)	(V_{15}, U_{15})	(24,11)	
(V_{16}, U_{16})	(32,21)	(V_{16}, U_{16})	(29,12)	
(V_{17}, U_{17})	(12,10)	(V_{17}, U_{17})	(3,7)	
(V_{18}, U_{18})	(13,0)	(V_{18}, U_{18})	(25,15)	
(V_{19}, U_{19})	(12,23)	(V_{19}, U_{19})	(18,32)	
(V_{20}, U_{20})	(32,12)	(V_{20}, U_{20})	(2,0)	
(V_{21}, U_{21})	(9,22)	(V_{21}, U_{21})	(18,1)	
(V_{22}, U_{22})	(10,21)	(V_{22}, U_{22})	(25,18)	
(V_{23}, U_{23})	(12,32)	(V_{23}, U_{23})	(3,26)	
(V_{24}, U_{24})	(2,0)	(V_{24}, U_{24})	(29,21)	
(V_{25}, U_{25})	(12,1)	(V_{25}, U_{25})	(24,22)	
(V_{26}, U_{26})	(10,12)	(V_{26}, U_{26})	(7,12)	
(V_{27}, U_{27})	(9,11)	(V_{27}, U_{27})	(3,29)	
(V_{28}, U_{28})	(32,21)	(V_{28}, U_{28})	(14,15)	
(V_{29}, U_{29})	(12,10)	(V_{29}, U_{29})	(18,10)	
(V_{30}, U_{30})	(13,0)	(V_{30}, U_{30})	(13,0)	
(V_{31}, U_{31})	(12,23)	(V_{31}, U_{31})	(18,23)	
(V_{32}, U_{32})	(32,12)	(V_{32}, U_{32})	(14,18)	
(V_{33}, U_{33})	(9,22)	(V_{33}, U_{33})	(3,4)	

TABLE 1. Values of $L(\Delta, 33)$ for $(\frac{\Delta}{33}) = 1$ and $(\frac{\Delta}{33}) = -1$

	$L(\Delta$,37)	
$a = 17, \Delta = 26, (\frac{\Delta}{N}) = 1 \text{ and } S(N) = 36$		$a = 11, \Delta = 6, (\frac{\Delta}{N}) = -1 \text{ and } S(N) = 38$	
(V_0, U_0)	(2,0)	(V_0, U_0)	(2,0)
(V_1, U_1)	(17,1)	(V_1, U_1)	(11,1)
(V_2, U_2)	(28,17)	(V_2,U_2)	(8,11)
(V_3, U_3)	(15,29)	(V_3, U_3)	(3,9)
(V_4,U_4)	(5,32)	(V_4, U_4)	(25,14)
(V_5, U_5)	(33,34)	(V_5, U_5)	(13,34)
(V_6, U_6)	(1,28)	(V_6, U_6)	(7,27)
(V_7, U_7)	(21,35)	(V_7, U_7)	(27,4)
(V_8,U_8)	(23,12)	(V_8, U_8)	(31,17)
(V_9,U_9)	(0,21)	(V_9,U_9)	(18,35)
(V_{10}, U_{10})	(14,12)	(V_{10}, U_{10})	(19,35)
(V_{11}, U_{11})	(16,35)	(V_{11}, U_{11})	(6,7)
(V_{12}, U_{12})	(36,28)	(V_{12}, U_{12})	(10,4)
(V_{13}, U_{13})	(4,34)	(V_{13}, U_{13})	(30,27)
(V_{14}, U_{14})	(32,32)	(V_{14}, U_{14})	(24,34)
(V_{15}, U_{15})	(22,29)	(V_{15}, U_{15})	(12,14)
(V_{16}, U_{16})	(9,27)	(V_{16}, U_{16})	(34,9)
(V_{17}, U_{17})	(20,1)	(V_{17}, U_{17})	(29,11)
(V_{18}, U_{18})	(35,0)	(V_{18}, U_{18})	(26,1)
(V_{19}, U_{19})	(20,36)	(V_{19}, U_{19})	(35,0)
(V_{20}, U_{20})	(9,3)	(V_{20}, U_{20})	(26,36)
(V_{21}, U_{21})	(22,8)	(V_{21}, U_{21})	(29,26)
(V_{22}, U_{22})	(32,5)	(V_{22}, U_{22})	(34,28)
(V_{23}, U_{23})	(4,3)	(V_{23}, U_{23})	(12,23)
(V_{24}, U_{24})	(36,9)	(V_{24}, U_{24})	(24,3)
(V_{25}, U_{25})	(16,2)	(V_{25}, U_{25})	(30,10)
(V_{26}, U_{26})	(14,25)	(V_{26}, U_{26})	(10,33)
(V_{27}, U_{27})	(0,16)	(V_{27}, U_{27})	(6,20)
(V_{28},U_{28})	(23,25)	(V_{28}, U_{28})	(19,2)
(V_{29}, U_{29})	(21,2)	(V_{29}, U_{29})	(18,2)
(V_{30}, U_{30})	(1,9)	(V_{30}, U_{30})	(31,20)
(V_{31}, U_{31})	(33,3)	(V_{31}, U_{31})	(27,33)
(V_{32}, U_{32})	(5,5)	(V_{32}, U_{32})	(7,10)
(V_{33}, U_{33})	(15,8)	(V_{33}, U_{33})	(13,3)
(V_{34}, U_{34})	(28,20)	(V_{34}, U_{34})	(25,23)
(V_{35}, U_{35})	(17,1)	(V_{35}, U_{35})	(3,28)
(V_{36}, U_{36})	(2,0)	(V_{36}, U_{36})	(8,26)
(V_{37}, U_{37})	(17,1)	(V_{37}, U_{37})	(11,36)
(V_{38}, U_{38})	(28,17)	(V_{38}, U_{38})	(2,0)

TABLE 2. Values of $L(\Delta, 37)$ for $(\frac{\Delta}{37}) = 1$ and $(\frac{\Delta}{37}) = -1$

In table 1 the shaded cells are depicting that the values of $\left(V_{N-\left(\frac{\Delta}{N}\right)}(a,1),U_{N-\left(\frac{\Delta}{N}\right)}(a,1)\right)$ are not equal to $\left(V_{0}(a,1),U_{0}(a,1)\right)$ for composite N=33, for all the choices of a and Δ and in table 2 the shaded cells are depicting that the values of $\left(V_{N-\left(\frac{\Delta}{N}\right)}(a,1),U_{N-\left(\frac{\Delta}{N}\right)}(a,1)\right)$ are equal to $\left(V_{0}(a,1),U_{0}(a,1)\right)$ for prime N=37, for all the choices of a and Δ .

4. Conclusion

There are several studies on Lucas sequences and their applications [7] [9]. Primality tests with Lucas sequences by Lucas and Lehmer given in [9] are based on factorization of $N \pm 1$. In this paper we proposed a primality test with pair of Lucas sequences $\left(V_n(a,1), U_n(a,1)\right)$ mod N from the set $L(\Delta,N)$ of S(N) distinct Lucas sequences for $S(N) = \text{LCM}\left[\left\{p_i^{e_i-1}\left(p_i-\left(\frac{\Delta}{p_i}\right)\right)\right\}_{i=1}^s\right]$. An algorithm for primality test given, employing the addition chain as in [11] for computation of $\left(V_n(a,1), U_n(a,1)\right)$.

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CONFLICT OF INTERESTS

The author(s) declare that there is no conflict of interests.

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